## Optimal multi-distribution learning



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"Optimal multi-distribution learning," Z. Zhang, W. Zhan, Y. Chen, S. Du, J. Lee, arXiv:2312.05134, 2023


In multi-distribution learning, an agent aims to learn a shared model to fit multiple (unknown) data distributions

- diverse data sources (e.g., localities, communities, populations)
- heterogeneous objectives $\longrightarrow$ need a balance

$\perp$ NewYork-Presbyterian
7 The University Hospital of Columbia and Cornell
$\longrightarrow \mathcal{D}_{3}$
- $k$ unknown data distributions $\mathcal{D}_{1}, \ldots, \mathcal{D}_{k}$ (e.g., localities, communities, populations)
- hypothesis class $\mathcal{H}$ : VC dimension $d$
- known loss function $\ell$ (e.g., misclassification error)

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goal: learn an $\varepsilon$-optimal $\underbrace{\text { hypothesis } \widehat{h}}_{\text {possibly random }}$ (in min-max sense)

$$
\max _{1 \leq i \leq k} \underset{(x, y) \sim \mathcal{D}_{i}, \widehat{h}}{\mathbb{E}}[\ell(\widehat{h},(x, y))] \leq \min _{h \in \mathcal{H}} \max _{1 \leq i \leq k} \underset{(x, y) \sim \mathcal{D}_{i}}{\mathbb{E}}[\ell(h,(x, y))]+\varepsilon
$$

Mohri et al. '19, Sagawa et al. '19, Blum et al. '17, Buhlmann et al. '15, Guo '23 ...

distributionally robust learning

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distributionally robust learning

collaborative learning

## Adaptive vs. non-adaptive sampling

- non-adaptive sampling: pre-determine sample-size budgets for each distribution beforehand
$\longrightarrow$ loss of data efficiency


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learning 1 distribution



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learning 1 distribution




## Prior works: VC classes

| paper | sample complexity |
| :---: | :---: |
| Haghtalab et al. '22 | $\frac{d+k}{\varepsilon^{2}}+\frac{d k}{\varepsilon}$ |
| Awasthi et al. '23 | $\frac{d}{\varepsilon^{4}}+\frac{k}{\varepsilon^{2}}$ |
| (lower bound) Haghtalab et al. '22 | $\frac{d+k}{\varepsilon^{2}}$ |

## Prior works: VC classes



# Can we close the gap between achievability and lower bound? 

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36 t Annual Conference on Learning Theory

Open Problem: The Sample Complexity of Multi-Distribution
Learning for VC Classes

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## Main results

## Theorem 1 (Zhang, Zhan, Chen, Du, Lee '23)

We can design an algorithm that returns randomized hypothesis $\widehat{h}$ s.t.

$$
\max _{1 \leq i \leq k} \underset{(x, y) \sim \mathcal{D}_{i}, \widehat{h}}{\mathbb{E}}[\ell(\widehat{h},(x, y))] \leq \min _{h \in \mathcal{H}} \max _{1 \leq i \leq k} \underset{(x, y) \sim \mathcal{D}_{i}}{\mathbb{E}}[\ell(h,(x, y))]+\varepsilon
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with sample complexity

$$
\widetilde{O}\left(\frac{d+k}{\varepsilon^{2}}\right)
$$

- matches the minimax lower bound (up to log factors)


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- solves a COLT open problem (concurrent work: Peng '23)
- can be extended to Rademacher classes
- algorithm is $\underbrace{\text { oracle-efficient }}$ (solves another COLT open problem) only needs to call ERM oracle


## Algorithm design



## A game-theoretic view


finding most favorable hypothesis

## Preliminaries: learning in games



- no-regret algorithm: online algorithm w/ $\underbrace{\text { sub-linear regret }}$ over any adversary

$$
\frac{1}{T} \operatorname{Regret}(T) \rightarrow 0
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- e.g., Hedge algorithm (equivalent to online mirror descent)


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- e.g., Hedge algorithm (equivalent to online mirror descent)
- best-response: play argmin or argmax (not always no-regret)

- min-player/max-player: no-regret/no-regret (Haghtalab et al. '22)

$$
\frac{d+k}{\varepsilon^{2}}+\frac{d k}{\varepsilon} \quad(\text { burn-in due to covering of } \mathcal{H})
$$

- min-player/max-player: best-response/no-regret (Awasthi et al. '23)

$$
\frac{d}{\varepsilon^{4}}+\frac{k}{\varepsilon^{2}} \quad \text { (lack of sample reuse) }
$$

## Our approach: best-response/no-regret

At iteration $t$ :

- min-player computes empirical best response

$$
\begin{gathered}
h^{t} \approx \arg \min _{h \in \mathcal{H}} L\left(h, w^{t}\right) \\
\circ L(h, w):=\sum_{i=1}^{k} w_{i} \underset{(x, y) \sim \mathcal{D}_{i}}{\mathbb{E}}[\ell(h,(x, y))] \text { (loss w.r.t. weighted dist) }
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w_{i}^{t} \propto w_{i}^{t-1} \exp \left(\eta \widehat{r}_{i}^{t}\right) \quad \text { with } \widehat{r}_{i}^{t}: \text { empirical risk for } \mathcal{D}_{i}
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$$

Output: randomized hypothesis $\widehat{h} \sim \operatorname{Uniform}\left(\left\{h^{t}\right\}_{1 \leq t \leq T}\right)$

## Key algorithmic distinction from prior work

## adaptive sampling $\quad+$ sample reuse <br> \#samples from $\mathcal{D}_{i}$ based on $\left\{w_{i}^{t}\right\}$

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Sampling strategy at iteration $t$ :

- best-response: $\underbrace{\text { have } \frac{d+k}{\varepsilon^{2}} w_{i}^{t} \text { samples available from } \mathcal{D}_{i}, ~}_{\text {reuse samples }}$


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Sampling strategy at iteration $t$ :

- best-response: have $\frac{d+k}{\varepsilon^{2}} \max _{1 \leq \tau \leq t} w_{i}^{\tau}$ samples available from $\mathcal{D}_{i}$ reuse samples


## Key algorithmic distinction from prior work



Sampling strategy at iteration $t$ :

- best-response: have $\frac{d+k}{\varepsilon^{2}} \max _{1 \leq \tau \leq t} w_{i}^{\tau}$ samples available from $\mathcal{D}_{i}$
reuse samples
- no-regret: draw $k \max _{1 \leq \tau \leq t} w_{i}^{\tau}$ samples from $\mathcal{D}_{i}$
fresh samples


## Key technical challenges

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$$

concentration + doubling trick + combinatorics

## Concurrent work: Peng et al. '23

Peng et al. '23 established a sample complexity of

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- optimal up to some sub-polynomial term
- a very different algorithm
- recursive structure to eliminate non-optimal hypotheses


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Question: is it possible to find an $\varepsilon$-optimal deterministic hypothesis $\mathrm{w} /$ the same sample complexity (another COLT open problem)?

Answer: No!

- finding an $\varepsilon$-optimal deterministic policy needs $\Omega\left(\frac{d k}{\varepsilon^{2}}\right)$ samples


## Summary: multi-distribution learning



- settles the sample complexity of MDL under on-demand sampling
- solves 3 COLT open problems posed by Awasthi et al. '23


## Concluding remarks

Advancing frontier of statistical learning requires integrated thinking of modern statistics, optimization \& game theory

online learning \& games

(high-dimensional) statistics
"Optimal multi-distribution learning," Z. Zhang, W. Zhan, Y. Chen, S. Du, J. Lee, arXiv:2312.05134, 2023

