

## Homework 1

Due date: Wednesday, Mar. 1, 2017 (at the beginning of class)

You are allowed to drop 1 subproblem without penalty. In addition, up to 1 bonus point will be awarded to each subproblem for clean, well-organized, and elegant solutions.

## 1. Mutual coherence (40 points)

Recall that for an arbitrary pair of orthonormal bases  $\Psi = [\psi_1, \dots, \psi_n] \in \mathbb{R}^{n \times n}$  and  $\Phi = [\phi_1, \dots, \phi_n] \in \mathbb{R}^{n \times n}$ , the mutual coherence  $\mu(\Psi, \Phi)$  of these two bases is defined by

$$\mu(\Psi, \Phi) = \max_{1 \leq i, j \leq n} |\psi_i^\top \phi_j| \quad (1)$$

(a) Show that

$$\frac{1}{\sqrt{n}} \leq \mu(\Psi, \Phi) \leq 1.$$

(b) Let  $\Psi = \mathbf{I}$ , and suppose that  $\Phi = [\phi_{i,j}]_{1 \leq i, j \leq n}$  is a Gaussian random matrix such that the  $\phi_{i,j}$ 's are i.i.d. random variables drawn from  $\phi_{i,j} \sim \mathcal{N}(0, 1/n)$ . Can you provide an upper estimate on  $\mu(\Psi, \Phi)$  as defined in (1)? Since  $\Phi$  is a random matrix, we expect your answer to be a function  $f(n)$  such that  $\mathbb{P}\{\mu(\Psi, \Phi) > f(n)\} \rightarrow 0$  as  $n$  scales.

Hint: to simplify analysis, you are allowed to use the crude approximation  $\mathbb{P}\{|z| > \tau\} \approx \exp(-\tau^2/2)$  for large  $\tau > 0$ , where  $z \sim \mathcal{N}(0, 1)$ .

(c) Set  $n = 100$ . Generate a random matrix  $\Phi$  as in Part (b), and compute  $\mu(\mathbf{I}, \Phi)$ . Report the empirical distribution (i.e. histogram) of  $\mu(\mathbf{I}, \Phi)$  out of 1000 simulations. How does your simulation result compare to your estimate in Part (b)?

(d) We now generalize the mutual coherence measure to accommodate a more general set of vectors beyond two bases. Specifically, for any given matrix  $\mathbf{A} = [\mathbf{a}_1, \dots, \mathbf{a}_p] \in \mathbb{R}^{n \times p}$  obeying  $n \leq p$ , define the mutual coherence of  $\mathbf{A}$  as

$$\mu(\mathbf{A}) = \max_{1 \leq i, j \leq p, i \neq j} \left| \frac{\mathbf{a}_i^\top \mathbf{a}_j}{\|\mathbf{a}_i\| \|\mathbf{a}_j\|} \right|.$$

Show that

$$\mu(\mathbf{A}) \geq \sqrt{\frac{p-n}{p-1} \cdot \frac{1}{n}}.$$

This is a special case of the Welch bound.

Hint: you may want to use the following inequality: for any positive semidefinite  $\mathbf{M} \in \mathbb{R}^{n \times n}$ ,  $\|\mathbf{M}\|_{\mathbb{F}}^2 \geq \frac{1}{n} (\sum_{i=1}^n \lambda_i(\mathbf{M}))^2$ .

## 2. Picket-fence signal (10 points)

Suppose that  $\sqrt{n}$  is an integer. Let  $\mathbf{x} \in \mathbb{R}^n$  be a picket-fence signal with uniform spacing  $\sqrt{n}$  such that

$$x_i = \begin{cases} 1, & \text{if } \frac{i-1}{\sqrt{n}} \text{ is an integer,} \\ 0, & \text{else,} \end{cases} \quad i = 1, \dots, n.$$

Compute

$$\|\mathbf{x}\|_0 \cdot \|\mathbf{F}\mathbf{x}\|_0 \quad \text{and} \quad \|\mathbf{x}\|_0 + \|\mathbf{F}\mathbf{x}\|_0,$$

where  $\mathbf{F}$  is the Fourier matrix such that

$$(\mathbf{F})_{k,l} = \frac{1}{\sqrt{n}} \exp\left(-i \frac{2\pi(k-1)(l-1)}{n}\right), \quad 1 \leq k, l \leq n.$$

How do they compare to the uncertainty principles we derive in class?

### 3. $\ell_1$ minimization (20 points)

Suppose that  $\mathbf{A}$  is an  $n \times 2n$  dimensional matrix. Let  $\mathbf{x} \in \mathbb{R}^{2n}$  be an unknown  $k$ -sparse vector, and  $\mathbf{y} = \mathbf{A}\mathbf{x}$  the observed system output. This problem is concerned with  $\ell_1$  minimization (or basis pursuit) in recovering  $\mathbf{x}$ , i.e.

$$\text{minimize}_{\mathbf{z} \in \mathbb{R}^{2n}} \|\mathbf{z}\|_1 \quad \text{s.t.} \quad \mathbf{A}\mathbf{z} = \mathbf{y}. \quad (2)$$

(a) An optimization problem is called a linear program (LP) if it has the form

$$\begin{aligned} \text{minimize}_{\mathbf{z}} \quad & \mathbf{c}^\top \mathbf{z} + \mathbf{d} \\ \text{s.t.} \quad & \mathbf{G}\mathbf{z} \leq \mathbf{h} \\ & \mathbf{A}\mathbf{z} = \mathbf{b} \end{aligned}$$

where  $\mathbf{c}, \mathbf{d}, \mathbf{G}, \mathbf{h}, \mathbf{A}$ , and  $\mathbf{b}$  are known. Here, for any two vectors  $\mathbf{r}$  and  $\mathbf{s}$ , we say  $\mathbf{r} \leq \mathbf{s}$  if  $r_i \leq s_i$  for all  $i$ . Show that (2) can be converted to a linear program.

(b) Set  $n = 256$ , and let  $k$  range between 1 and 128. For each choice of  $k$ , run 10 independent numerical experiments: in each experiment, generate  $\mathbf{A} = [a_{i,j}]_{1 \leq i \leq n, 1 \leq j \leq 2n}$  as a random matrix such that the  $a_{i,j}$ 's are i.i.d. standard Gaussian random variables, generate  $\mathbf{x} \in \mathbb{R}^{2n}$  as a random  $k$ -sparse signal (e.g. you may generate the support of  $\mathbf{x}$  uniformly at random, with each non-zero entry drawn from the standard Gaussian distribution), and solve (2) with  $\mathbf{y} = \mathbf{A}\mathbf{x}$ . An experiment is claimed successful if the solution  $\mathbf{z}$  returned by (2) obeys  $\|\mathbf{x} - \mathbf{z}\|_2 \leq 0.001\|\mathbf{x}\|_2$ . Report the empirical success rates (averaged over 10 experiments) for each choice of  $k$ .